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Distributive laws in residuated binars

Wesley Fussner and Peter Jipsen

Abstract. In residuated binars there are six non-obvious distributivity identities of $\cdot, /, \backslash$ over \wedge, \vee . We show that in residuated binars with distributive lattice reducts there are some dependencies among these identities; specifically, there are six pairs of identities that imply another one of these identities, and we provide counterexamples to show that no other dependencies exist among these.

Mathematics Subject Classification. 06F05, 03G10, 08B15.

Keywords. Residuated lattices, residuated binars, residuation, subvariety lattices.

1. Introduction

A *residuated binar* is an algebra $\mathbf{A} = (A, \wedge, \vee, \cdot, \backslash, /)$, where (A, \wedge, \vee) is a lattice, \cdot is a binary operation on A , and for all $x, y, z \in A$,

$$x \cdot y \leq z \iff x \leq z/y \iff y \leq x \backslash z.$$

A *residuated semigroup* is a residuated binar for which \cdot is associative, and a residuated binar possessing an identity element e for \cdot is called *unital*. An expansion of a unital residuated semigroup by a constant designating the identity is called a *residuated lattice* [5]. Note that to reduce the need for parentheses, we assume that \cdot has priority over $\backslash, /$, which in turn have priority over \wedge, \vee . We also write $x \cdot y$ as xy . Using these conventions, all of the aforementioned algebras satisfy the distributive laws

$$x(y \vee z) = xy \vee xz \tag{ $\cdot \vee$ }$$

$$(x \vee y)z = xz \vee yz \tag{ $\vee \cdot$ }$$

$$x \backslash (y \wedge z) = x \backslash y \wedge x \backslash z \tag{ $\backslash \wedge$ }$$

$$(x \wedge y)/z = x/z \wedge y/z \tag{ $\wedge /$ }$$

$$x/(y \vee z) = x/y \wedge x/z \tag{ $/ \vee$ }$$

$$(x \vee y) \backslash z = x \backslash z \wedge y \backslash z \tag{ $\vee \backslash$ }$$

However, in general neither lattice distributivity nor any of the equations

$$\begin{aligned}
x(y \wedge z) &= xy \wedge xz && (\cdot \wedge) \\
(x \wedge y)z &= xz \wedge yz && (\wedge \cdot) \\
x \setminus (y \vee z) &= x \setminus y \vee x \setminus z && (\setminus \vee) \\
(x \vee y) / z &= x / z \vee y / z && (\vee /) \\
(x \wedge y) \setminus z &= x \setminus z \vee y \setminus z && (\wedge \setminus) \\
x / (y \wedge z) &= x / y \vee x / z && (/ \wedge)
\end{aligned}$$

hold in these algebras.

If t is a term in the language of residuated binars (or residuated semigroups), then the *opposite of t* is the term t^{op} defined recursively as follows. For x a variable, set $x^{\text{op}} = x$, and if s and t are terms then set $(s \cdot t)^{\text{op}} = t^{\text{op}} \cdot s^{\text{op}}$, $(s/t)^{\text{op}} = t^{\text{op}} \setminus s^{\text{op}}$, $(s \setminus t)^{\text{op}} = t^{\text{op}} / s^{\text{op}}$, $(s \wedge t)^{\text{op}} = t^{\text{op}} \wedge s^{\text{op}}$, and $(s \vee t)^{\text{op}} = t^{\text{op}} \vee s^{\text{op}}$ (and $e^{\text{op}} = e$ in the presence of a multiplicative identity e). The opposite of an equation $s = t$ is defined by $(s = t)^{\text{op}} = (s^{\text{op}} = t^{\text{op}})$. *Mirror duality* for residuated binars provides that an equation ε holds in the variety of all residuated binars if and only if ε^{op} does as well. If $\Sigma \cup \{\varepsilon\}$ is a set of equations in the language of residuated binars and $\Sigma^{\text{op}} = \{\sigma^{\text{op}} : \sigma \in \Sigma\}$, then $\Sigma \models \varepsilon$ holds in the variety of residuated binars if and only if $\Sigma^{\text{op}} \models \varepsilon^{\text{op}}$ holds. Observe that $(\cdot \wedge)^{\text{op}}$, $(\setminus \vee)^{\text{op}}$, and $(\wedge \setminus)^{\text{op}}$ are respectively $(\wedge \cdot)$, $(\vee /)$, and $(/ \wedge)$.

In the presence of a multiplicative identity e , left and right prelinearity

$$e \leq x \setminus y \vee y \setminus x \tag{lp}$$

$$e \leq x / y \vee y / x, \tag{rp}$$

have a connection to the six nontrivial distributive laws given above. In particular, [2, Proposition 6.10] shows that in residuated lattices satisfying e -distributivity

$$(x \vee y) \wedge e = (x \wedge e) \vee (y \wedge e), \tag{ed}$$

the equations (lp) , $(\wedge \setminus)$, and $(\setminus \vee)$ are pairwise equivalent, as are the equations (rp) , $(/ \wedge)$, and $(\vee /)$. Because (lp) and (rp) axiomatize semilinear residuated lattices (i.e., those that are subdirect products of totally-ordered residuated lattices) under appropriate technical hypotheses (see [2]), this provides one explanation of the well-known fact that all six nontrivial distributive laws hold in semilinear residuated lattices. However, a residuated lattice may satisfy all six nontrivial distributive laws even though it is not semilinear (this is the case, e.g., in lattice-ordered groups).

The dependencies among the six nontrivial distributive laws are more complicated in the absence of a multiplicative identity. Sections 2 and 3 provide a complete description of the dependencies among the nontrivial distributive laws under the hypothesis of lattice distributivity, both for residuated binars and residuated semigroups. Section 4 provides some additional implications among the distributive laws in unital residuated binars, and in the presence of lattice complements. We conclude in Section 5 by proposing some open problems.

2. Implications among the nontrivial distributive laws

A residuated binar with a distributive lattice reduct may be associated with its *frame*. The frame of a lattice-distributive residuated binar \mathbf{A} may be obtained by taking the set of prime filters of the lattice reduct of \mathbf{A} , and endowing it with the order \leq given by reverse subset inclusion as well as the ternary relation R defined by

$$R(F, G, H) \iff F \leq G \cdot H,$$

where $G \cdot H$ is the filter generated by the complex product $\{xy : x \in G, y \in H\}$ of G and H . Note that with respect to \leq , R is antitone in its first coordinate and isotone in its second and third coordinates.

Satisfaction of either of the identities $(\setminus \vee)$ and $(\vee /)$ has significant consequences for the frame of a lattice-distributive residuated binar [4], and the nontrivial distributive laws may be profitably analyzed from the point of view of frames. In fact, for lattice-distributive residuated binars, each of the distributive laws introduced in the previous section may be rendered in terms of an equivalent first-order condition on the corresponding frames by application of the Ackermann-lemma-based algorithm (ALBA) [3]. The elements of the frame of an algebra correspond to completely join-irreducible elements (= principal downsets) in the algebra of downsets of the frame. Moreover, such algebras have complete distributive lattice reducts in which all elements are joins of completely join-irreducibles and meets of completely meet-irreducibles. We illustrate the ALBA algorithm for such algebras by deriving the equivalent first-order condition for $(\vee /)$, i.e., $(x \vee y)/z = x/z \vee y/z$. This identity holds if and only if for all completely join-irreducibles j we have that $j \leq (x \vee y)/z \implies j \leq x/z \vee y/z$. Since join-irreducibles in distributive lattices are join-prime, this holds if and only if

$$jz \leq x \vee y \implies jz \leq x \text{ or } jz \leq y,$$

or, equivalently, for all completely join-irreducibles u, v, w

$$(\forall w(w \leq jz \implies w \leq x \vee y)) \implies (u \leq jz \implies u \leq x) \text{ or } (v \leq jz \implies v \leq y).$$

Since $u \leq jz$ is equivalent to $\exists p(p \leq z \ \& \ R(u, j, p))$ and $v \leq jz$ is equivalent to $\exists q(q \leq z \ \& \ R(v, j, q))$, we instantiate $z = p \vee q$ to obtain for all completely join-irreducibles j, u, v, p, q ,

$$\begin{aligned} & (\forall w(w \leq j(p \vee q) \implies w \leq x \vee y)) \\ & \implies (R(u, j, p) \implies u \leq x) \text{ or } (R(v, j, q) \implies v \leq y). \end{aligned}$$

In a complete distributive lattice, for any completely join-irreducible u there exists a corresponding completely meet-irreducible u^* such that $u \not\leq x$ if and only if $x \leq u^*$. Instantiating x, y by u^*, v^* produces

$$(\forall w(w \leq j(p \vee q) \implies w \leq u^* \vee v^*)) \implies (\neg R(u, j, p)) \text{ or } (\neg R(v, j, q)).$$

Since w is join-irreducible, $w \leq j(p \vee q)$ is equivalent to $R(w, j, p)$ or $R(w, j, q)$, and $w \leq u^* \vee v^*$ is equivalent to $u \not\leq w$ or $v \not\leq w$. Rearranging the formula

contrapositively then produces

$$[R(u, j, p) \ \& \ R(v, j, q)] \implies \exists w[u, v \leq w \ \& \ (R(w, j, p) \ \text{or} \ R(w, j, q))].$$

Conversely, suppose that a frame satisfies this condition. To see that $(\vee/)$ holds in the algebra of downsets, let x, y, z be downsets of the frame and suppose that $j \in (x \cup y)/z$. Then $\downarrow j \cdot z \subseteq x \cup y$. We assume that $\downarrow j \not\subseteq x/z$ (i.e., that $\downarrow j \cdot z \not\subseteq x$), and we aim to show that $\downarrow j \cdot z \subseteq y$ (i.e., $j \in y/z$). Note that $\downarrow j \cdot z \not\subseteq x$ implies that there exists $u \in \downarrow j \cdot z$ such that $u \notin x$, and this gives that there exists $p \in z$ such that $R(u, j, p)$. Let $v \in \downarrow j \cdot z$. Then there exists $q \in z$ such that $R(v, j, q)$. Applying the hypothesis to $R(u, j, p)$ and $R(v, j, q)$ provides that there exists w with $u, v \leq w$ and $R(w, j, p)$ or $R(w, j, q)$. If $R(w, j, p)$ holds, then $w \in \downarrow j \cdot z \subseteq x \cup y$, whence $w \in x$ or $w \in y$. If $w \in x$, then $u \leq w$ and x being a downset imply that $u \in x$, a contradiction. Hence $w \in y$, and since $v \leq w$ and y is a downset we get $v \in y$. Hence $\downarrow j \cdot z \subseteq y$, proving the result.

Similarly $(\wedge \setminus)$ is equivalent to the condition that for all j, u, v, p, q ,

$$[R(p, u, j) \ \& \ R(q, v, j)] \implies \exists w[w \leq u, v \ \& \ (R(p, w, j) \ \text{or} \ R(q, w, j))],$$

whereas $(\setminus \vee)$ is equivalent to the condition that for all j, u, v, p, q ,

$$[R(u, p, j) \ \& \ R(v, q, j)] \implies \exists w[u, v \leq w \ \& \ (R(w, p, j) \ \text{or} \ R(w, q, j))].$$

Proposition 2.1. *Let \mathbf{A} be a residuated binar with a distributive lattice reduct. If \mathbf{A} satisfies both $(\vee/)$ and $(\wedge \setminus)$, then \mathbf{A} also satisfies $(\setminus \vee)$.*

Proof. Suppose that both $(\vee/)$ and $(\wedge \setminus)$ hold. We use the equivalent frame conditions to verify $(\setminus \vee)$, so suppose that j, u, v, p, q are points in the frame of \mathbf{A} such that $R(u, p, j)$ and $R(v, q, j)$. By the frame condition for $(\wedge \setminus)$ there exists w' with $w' \leq p, q$ and one of $R(u, w', j)$ or $R(v, w', j)$. Suppose first that $R(u, w', j)$ holds. Then $R(u, w', j)$ and $R(v, q, j)$, and by monotonicity and $w' \leq q$ we have $R(u, q, j)$ and $R(v, q, j)$. Using the frame condition for $(\vee/)$ we obtain w such that $u, v \leq w$ and $R(w, q, j)$. On the other hand, if $R(v, w', j)$ holds then $R(v, w', j)$ and $R(u, p, j)$. Monotonicity and $w' \leq p$ then gives $R(v, p, j)$ and $R(u, p, j)$, and by the frame condition for $(\vee/)$ there exists w with $u, v \leq w$ and $R(w, p, j)$. In either case, there exists w with $u, v \leq w$ and either $R(w, p, j)$ or $R(w, q, j)$, which completes the proof. \square

Other results of this kind may be discovered by appealing to equivalent conditions on frames. However, an entirely algebraic treatment is also possible. The next lemma is an important step in this direction.

Lemma 2.2. *Each of the following gives a pair of identities that are equivalent in residuated binars.*

- (1) $(\cdot \wedge)$ and $xz \wedge yw \leq (x \vee y)(z \wedge w)$.
- (2) $(\wedge \cdot)$ and $xz \wedge yw \leq (x \wedge y)(z \vee w)$.
- (3) $(\setminus \vee)$ and $(x \vee y) \setminus (z \vee w) \leq x \setminus z \vee y \setminus w$.
- (4) $(\vee/)$ and $(z \vee w) / (x \vee y) \leq z / x \vee w / y$.
- (5) $(\wedge \setminus)$ and $(x \wedge y) \setminus (z \wedge w) \leq x \setminus z \vee y \setminus w$.
- (6) $(/ \wedge)$ and $(z \wedge w) / (x \wedge y) \leq z / x \vee w / y$.

Proof. We prove (1) and (3); (2) and (4) follow by a symmetric argument, and (5) and (6) follow by a proof similar to (3) and (4).

For (1), note that if $xz \wedge yw \leq (x \vee y)(z \wedge w)$ holds then by instantiating $y = x$ we obtain $xz \wedge xw \leq x(z \wedge w)$. The reverse inequality follows from the isotonicity of multiplication, so $(\cdot \wedge)$ holds. Conversely, if $(\cdot \wedge)$ holds then we have $xz \wedge yw \leq (x \vee y)z \wedge (x \vee y)w = (x \vee y)(z \wedge w)$.

For (3), taking $y = x$ in the inequality $(x \vee y) \setminus (z \vee w) \leq x \setminus z \vee y \setminus w$ gives $x \setminus (z \vee w) \leq x \setminus z \vee x \setminus w$. The reverse inequality holds because \setminus is isotone in its numerator, whence $(\setminus \vee)$ holds. For the converse, note that $(\setminus \vee)$ implies $(x \vee y) \setminus (z \vee w) = (x \vee y) \setminus z \vee (x \vee y) \setminus w \leq x \setminus z \vee y \setminus w$, where the last step follows because \setminus is antitone in its denominator. \square

Theorem 2.3. *Let \mathbf{A} be a residuated binar with a distributive lattice reduct. Then:*

- (1) *If \mathbf{A} satisfies both $(\vee /)$ and $(\wedge \setminus)$, then \mathbf{A} also satisfies $(\setminus \vee)$.*
- (2) *If \mathbf{A} satisfies both $(\setminus \vee)$ and $(/\wedge)$, then \mathbf{A} also satisfies $(\vee /)$.*
- (3) *If \mathbf{A} satisfies both $(\cdot \wedge)$ and $(\vee /)$, then \mathbf{A} also satisfies $(/\wedge)$.*
- (4) *If \mathbf{A} satisfies both $(\wedge \cdot)$ and $(\setminus \vee)$, then \mathbf{A} also satisfies $(\wedge \setminus)$.*
- (5) *If \mathbf{A} satisfies both $(\wedge \setminus)$ and $(\cdot \wedge)$, then \mathbf{A} also satisfies $(\wedge \cdot)$.*
- (6) *If \mathbf{A} satisfies both $(/\wedge)$ and $(\wedge \cdot)$, then \mathbf{A} also satisfies $(\cdot \wedge)$.*

Proof. We provide proofs for (1) and (5); (2) and (6) follow by mirror duality. The others follow similarly.

For (1), suppose that $u \leq (x \vee y) \setminus (z \vee w)$. Then by residuation we get $x, y \leq x \vee y \leq (z \vee w) / u$, and by $(\vee /)$ we have $x \leq z / u \vee w / u$ and also $y \leq z / u \vee w / u$. Observe that $x = x \wedge (z / u \vee w / u)$ and $y = y \wedge (z / u \vee w / u)$, and by distributivity we obtain that $x = x_1 \vee x_2$ and $y = y_1 \vee y_2$, where

$$x_1 = x \wedge (z / u),$$

$$x_2 = x \wedge (w / u),$$

$$y_1 = y \wedge (z / u),$$

$$y_2 = y \wedge (w / u).$$

Note that

$$x_1 \leq z / u \implies u \leq x_1 \setminus z \leq (x_1 \wedge y_2) \setminus z,$$

$$x_2 \leq w / u \implies u \leq x_2 \setminus w \leq (x_2 \wedge y_1) \setminus w,$$

$$y_1 \leq z / u \implies u \leq y_1 \setminus z \leq (x_2 \wedge y_1) \setminus z,$$

$$y_2 \leq w / u \implies u \leq y_2 \setminus w \leq (x_1 \wedge y_2) \setminus w.$$

Hence we get that $u \leq (x_1 \wedge y_2) \setminus (z \wedge w) \leq x_1 \setminus z \vee y_2 \setminus w$ and likewise $u \leq (x_2 \wedge y_1) \setminus (z \wedge w) \leq x_2 \setminus z \vee y_1 \setminus w$. Also, $u \leq x_1 \setminus z \leq x_1 \setminus z \vee y_1 \setminus w$ and

$u \leq y_2 \setminus w \leq x_2 \setminus z \vee y_2 \setminus w$. This implies that:

$$\begin{aligned}
 u &\leq (x_1 \setminus z \vee y_2 \setminus w) \wedge (x_2 \setminus z \vee y_1 \setminus w) \wedge (x_1 \setminus z \vee y_1 \setminus w) \wedge (x_2 \setminus z \vee y_2 \setminus w) \\
 &= ((x_2 \setminus z \wedge x_1 \setminus z) \vee y_1 \setminus w) \wedge ((x_1 \setminus z \wedge x_2 \setminus z) \vee y_2 \setminus w) \\
 &= (x_1 \setminus z \wedge x_2 \setminus z) \vee (y_1 \setminus w \wedge y_2 \setminus w) \\
 &= (x_1 \vee x_2) \setminus z \vee (y_1 \vee y_2) \setminus w \\
 &= x \setminus z \vee y \setminus w.
 \end{aligned}$$

This proves that $(x \vee y) \setminus (z \vee w) \leq x \setminus z \vee y \setminus w$, whence (1) follows by Lemma 2.2(3).

To prove (5), suppose that $(x \wedge y)(z \vee w) \leq u$. By residuating and $(\wedge \setminus)$, we obtain $z, w \leq z \vee w \leq (x \wedge y) \setminus u = x \setminus u \vee y \setminus u$. Define

$$\begin{aligned}
 z_1 &= z \wedge (x \setminus u), \\
 z_2 &= z \wedge (y \setminus u), \\
 w_1 &= w \wedge (x \setminus u), \\
 w_2 &= w \wedge (y \setminus u),
 \end{aligned}$$

and note that by the distributivity of the lattice reduct we have $z = z_1 \vee z_2$ and $w = w_1 \vee w_2$. This provides

$$\begin{aligned}
 z_1 &\leq x \setminus u \implies xz_1 \leq u, \\
 z_2 &\leq y \setminus u \implies yz_2 \leq u, \\
 w_1 &\leq x \setminus u \implies xw_1 \leq u, \\
 w_2 &\leq y \setminus u \implies yw_2 \leq u,
 \end{aligned}$$

whence from the isotonicity of multiplication and the middle two items above, we obtain that $y(z_2 \wedge w_1) \leq u$ and $x(z_2 \wedge w_1) \leq u$. This provides that $(x \vee y)(z_2 \wedge w_1) = x(z_2 \wedge w_1) \vee y(z_2 \wedge w_1) \leq u$, and from the assumption $(\cdot \wedge)$ and Lemma 2.2(1) we conclude that $xz_2 \wedge yw_1 \leq u$. Now note that

$$\begin{aligned}
 xz \wedge yw &= x(z_1 \vee z_2) \wedge y(w_1 \vee w_2) \\
 &= (xz_1 \vee xz_2) \wedge (yw_1 \vee yw_2) \\
 &= (xz_1 \wedge yw_1) \vee (xz_1 \wedge yw_2) \vee (xz_2 \wedge yw_1) \vee (xz_2 \wedge yw_2) \\
 &\leq u,
 \end{aligned}$$

where the third equation above follows from lattice distributivity. It follows that $xz \wedge yw \leq (x \wedge y)(z \vee w)$, so $(\wedge \cdot)$ follows by Lemma 2.2(2). This gives (5). \square

The implications articulated in Theorem 2.3 are described by the directed graph in Figure 1. Each pair of identities given on the left-hand side (respectively, right-hand side) of the graph jointly imply their common successor on the right-hand side (respectively, left-hand side). Note that these consequences are hidden in the special case of e -distributive residuated lattices addressed in [2], where taken individually $(\wedge \setminus)$ and $(\setminus \vee)$ are equivalent, as are $(\vee /)$ and $(/ \wedge)$.

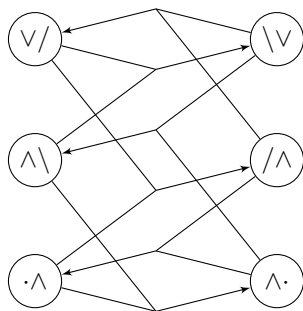


FIGURE 1. Dependencies among the nontrivial distributive laws.

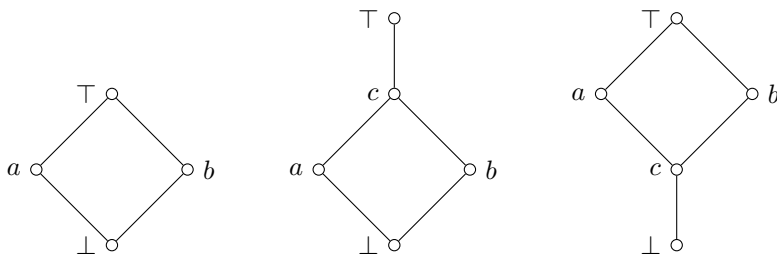


FIGURE 2. Labeled Hasse diagrams for the lattice reducts of $\mathbf{A}_1, \mathbf{A}_2, \mathbf{A}_3$ (left), $\mathbf{A}_4, \mathbf{A}_5$ (middle) and \mathbf{A}_6 (right).

3. The poset of subvarieties

The class of residuated binars with distributive lattice reducts forms a finitely-based variety \mathbf{RB} , and the implications announced in Theorem 2.3 entail some inclusions among the subvarieties of \mathbf{RB} determined by the nontrivial distributive laws. We will show that these are all of the inclusions among such subvarieties, completely describing the subposet of the subvariety lattice of \mathbf{RB} whose elements are axiomatized (modulo the theory of \mathbf{RB}) by any collection of the nontrivial distributive laws. The same analysis holds for residuated semigroups as well.

Proposition 3.1. *Theorem 2.3 gives the only implications among the six nontrivial distributive laws modulo the theory of residuated binars. The same holds for residuated semigroups.*

Proof. For each $i \in \{1, 2, 3, 4, 5, 6\}$ we define a residuated binar \mathbf{A}_i . The lattice reduct of each \mathbf{A}_i is given in Figure 2. We provide operation tables for \cdot in each \mathbf{A}_i below; the operation tables for \setminus and $/$ are uniquely determined

by these in each case. For \mathbf{A}_1 , \mathbf{A}_2 , and \mathbf{A}_3 :

\cdot	\perp	a	b	\top
\perp	\perp	\perp	\perp	\perp
a	\perp	\perp	\perp	\perp
b	\perp	\perp	\top	\top
\top	\perp	\perp	\top	\top

\cdot	\perp	a	b	\top
\perp	\perp	\perp	\perp	\perp
a	\perp	\perp	\perp	\perp
b	\perp	a	b	\top
\top	\perp	a	b	\top

\cdot	\perp	a	b	\top
\perp	\perp	\perp	\perp	\perp
a	\perp	\perp	a	a
b	\perp	\perp	b	b
\top	\perp	\perp	\top	\top

For \mathbf{A}_4 , \mathbf{A}_5 , and \mathbf{A}_6 :

\cdot	\perp	a	b	c	\top
\perp	\perp	\perp	\perp	\perp	\perp
a	\perp	\top	\perp	\top	\top
b	\perp	b	\perp	b	b
c	\perp	\top	\perp	\top	\top
\top	\perp	\top	\perp	\top	\top

\cdot	\perp	a	b	c	\top
\perp	\perp	\perp	\perp	\perp	\perp
a	\perp	\top	b	\top	\top
b	\perp	\perp	\perp	\perp	\perp
c	\perp	\top	b	\top	\top
\top	\perp	\top	b	\top	\top

\cdot	\perp	a	b	c	\top
\perp	\perp	\perp	\perp	\perp	\perp
a	\perp	a	\perp	\perp	a
b	\perp	\perp	b	\perp	b
c	\perp	\perp	\perp	\perp	\perp
\top	\perp	a	b	\perp	\top

Direct calculation verifies that:

- $\mathbf{A}_1 \models (\wedge/\wedge), (\wedge\backslash), (\wedge\cdot), (\cdot\wedge)$ and $\mathbf{A}_1 \not\models (\backslash\vee), (\vee/)$.
- $\mathbf{A}_2 \models (\backslash\vee), (\wedge\backslash), (\wedge\cdot), (\cdot\wedge)$ and $\mathbf{A}_2 \not\models (\vee/), (\wedge/\wedge)$.
- $\mathbf{A}_3 \models (\vee/), (\wedge/\wedge), (\wedge\cdot), (\cdot\wedge)$ and $\mathbf{A}_3 \not\models (\backslash\vee), (\wedge\backslash)$.
- $\mathbf{A}_4 \models (\vee/), (\backslash\vee), (\wedge/\wedge), (\cdot\wedge)$ and $\mathbf{A}_4 \not\models (\wedge\backslash), (\wedge\cdot)$.
- $\mathbf{A}_5 \models (\vee/), (\backslash\vee), (\wedge\backslash), (\wedge\cdot)$ and $\mathbf{A}_5 \not\models (\wedge/\wedge), (\cdot\wedge)$.
- $\mathbf{A}_6 \models (\vee/), (\backslash\vee), (\wedge/\wedge), (\wedge\backslash)$ and $\mathbf{A}_6 \not\models (\cdot\wedge), (\wedge\cdot)$.

Let $\varepsilon \in \{(\vee/), (\backslash\vee), (\wedge/\wedge), (\wedge\backslash), (\wedge\cdot), (\cdot\wedge)\}$. Then there exists a unique implication listed in Theorem 2.3 having ε as its consequent. Let $\varepsilon_1, \varepsilon_2$ be the identities in the antecedent of the aforementioned implication. Then the above countermodels show that if $\varepsilon \notin \Sigma \subseteq \{(\vee/), (\backslash\vee), (\wedge/\wedge), (\wedge\backslash), (\wedge\cdot), (\cdot\wedge)\}$ and $\varepsilon_1 \notin \Sigma$ or $\varepsilon_2 \notin \Sigma$, then ε is not entailed by Σ .

Note that each \mathbf{A}_i , $i \in \{1, 2, 3, 4, 5, 6\}$, is an associative residuated binar. The result therefore holds for residuated semigroups as well. \square

The left-hand side of Figure 3 gives the Hasse diagram of the poset of subvarieties of RB determined by the six nontrivial distributive laws. The coatoms in this diagram are subvarieties axiomatized modulo RB by a single nontrivial distributive law, and the atoms are subvarieties axiomatized by one of the four-element subsets of $\{(\vee/), (\backslash\vee), (\wedge/\wedge), (\wedge\backslash), (\wedge\cdot), (\cdot\wedge)\}$ satisfied in one of the models \mathbf{A}_i given in the proof of Proposition 3.1. The meets in this diagram correspond to intersection of subvarieties, but in general the joins do not correspond to joins in the lattice of subvarieties. The same diagram describes the corresponding subvariety poset for residuated semigroups since the models \mathbf{A}_i , $i \in \{1, 2, 3, 4, 5, 6\}$, are associative.

When \cdot is commutative in a residuated binar \mathbf{A} , the two residuals satisfy $x \backslash y = y / x$ for all $x, y \in A$ and therefore \backslash and $/$ coincide. In this event, $(\backslash\vee)$ is equivalent to $(\vee/)$, $(\wedge\backslash)$ is equivalent to (\wedge/\wedge) , and $(\cdot\wedge)$ is equivalent to $(\wedge\cdot)$. The poset of subvarieties axiomatized by the three pairwise independent nontrivial distributive laws is pictured on the right-hand side of Figure 3. The correctness of this diagram can be verified by observing that the models \mathbf{A}_1

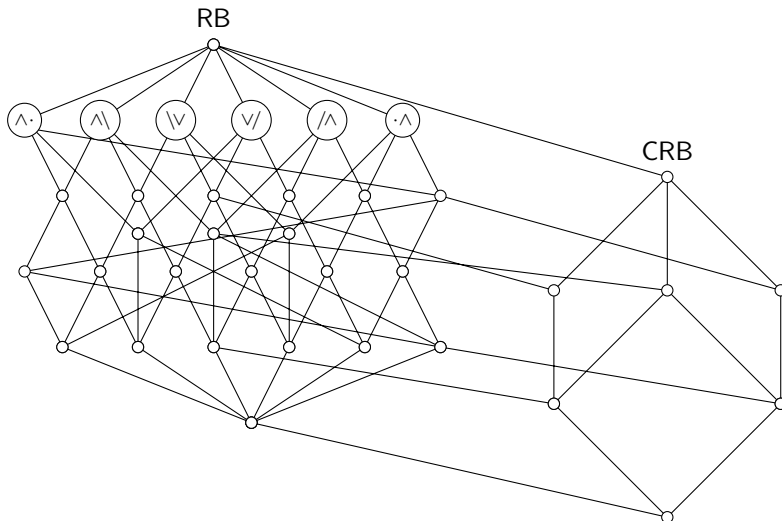


FIGURE 3. The poset of subvarieties determined by the non-trivial distributive laws in varieties of residuated binars RB and commutative residuated binars CRB.

and \mathbf{A}_6 are commutative. Since they are also associative, the same diagram describes the subvariety poset for commutative residuated semigroups.

4. Identity elements, complements, and prelinearity

We say that a residuated binar is *complemented* if its lattice reduct is complemented, and *Boolean* if its lattice reduct is a Boolean lattice. A unital residuated binar is called *integral* if it satisfies the identity $x \leq e$, where e is the multiplicative identity. We caution that this usage of *integral* is typical in the study of residuated lattices, but conflicts with the common usage in the theory of relation algebras. Boolean (unital) residuated binars are called (*u*)*r*-algebras in [6]. Note that if \cdot and \wedge coincide in a residuated binar \mathbf{A} , then \mathbf{A} is term-equivalent to a Brouwerian algebra (i.e., to the bottom-free reduct of a Heyting algebra). If additionally \mathbf{A} is a Boolean residuated binar, then \mathbf{A} is (term-equivalent to) a Boolean algebra.

The presence of complements and an identity element in a residuated binar can have a profound impact on whether it satisfies any of the six non-trivial distributive laws, a stark example of which is illustrated by the following lemma.

Lemma 4.1. *Let \mathbf{A} be a unital complemented residuated binar. If \mathbf{A} is integral, then \wedge and \cdot coincide.*

Proof. Since \mathbf{A} is integral, we have $x \cdot y \leq x \wedge y$ for all $x, y \in A$. This implies for any $x \in A$ we have that $x \cdot x' \leq x \wedge x' = \perp$, where x' is a complement of x .

On the other hand, since the identity element e is the greatest element of \mathbf{A} we have also that $x \vee x' = e$ for any $x \in A$. Multiplying by x and using $(\cdot \vee)$, we obtain $x = x \cdot e = x \cdot (x \vee x') = x^2 \vee x \cdot x' = x^2 \vee \perp = x^2$. This gives that \mathbf{A} is idempotent, whence for any $x, y \in A$, $x \wedge y = (x \wedge y) \cdot (x \wedge y) \leq x \cdot y \leq x \wedge y$, i.e., $x \cdot y = x \wedge y$. \square

Thus the only complemented integral residuated binars are Boolean algebras, which satisfy all six nontrivial distributive laws as well as lattice distributivity. Satisfaction of nontrivial distributive laws also often forces integrality in this setting.

Lemma 4.2. *Let \mathbf{A} be a unital residuated binar. If e has a complement e' and \mathbf{A} satisfies any one of the distributive laws $(\cdot \wedge)$, $(\wedge \cdot)$, $(\wedge \setminus)$, $(/\wedge)$, then \mathbf{A} is integral.*

Proof. We prove the result for $(\cdot \wedge)$ and $(\wedge \setminus)$. The result follows for $(\wedge \cdot)$ and $(/\wedge)$ by a symmetric argument.

First, suppose that \mathbf{A} satisfies $(\cdot \wedge)$. Then:

$$\begin{aligned} e' &= e \cdot e' \\ &\leq \top \cdot e' \\ &= \top \cdot e' \wedge \top \\ &= \top \cdot (e' \wedge e) \\ &= \top \cdot \perp \\ &= \perp \end{aligned}$$

where the last equality uses the identity $x \cdot \perp = \perp$, which holds in all residuated binars. It follows that $e = e \vee \perp = e \vee e' = \top$, hence $e = \top$.

Second, suppose that \mathbf{A} satisfies $(\wedge \setminus)$. Note that:

$$\begin{aligned} \top &= \perp \setminus \perp \\ &= (e \wedge e') \setminus \perp \\ &= (e \setminus \perp) \vee (e' \setminus \perp) \\ &= \perp \vee (e' \setminus \perp) \\ &= e' \setminus \perp, \end{aligned}$$

giving $\top \leq e' \setminus \perp$, and by residuation $e' \cdot \top \leq \perp$. As $e \leq \top$ and \cdot is isotone, we get $e' \cdot e \leq e' \cdot \top \leq \perp$. Therefore $e' \leq \perp$, so $e' = \perp$. It follows that $e = \top$, completing the proof. \square

Combining the previous two lemmas gives the following result.

Corollary 4.3. *Let \mathbf{A} be a complemented unital residuated binar. If \mathbf{A} satisfies any one of the distributive laws $(\cdot \wedge)$, $(\wedge \cdot)$, $(\wedge \setminus)$, $(/\wedge)$, then \mathbf{A} is a Boolean algebra.*

Proof. Since \mathbf{A} is complemented, e has a complement. Lemma 4.2 then gives that \mathbf{A} is integral, and so by Lemma 4.1 it follows that \mathbf{A} is a Boolean algebra. \square

Lemma 4.4. *Let \mathbf{A} be a unital Boolean residuated binar. If \mathbf{A} satisfies any one of the distributive laws $(\cdot\wedge)$, $(\wedge\cdot)$, $(\backslash\vee)$, $(\vee/)$, $(\wedge\backslash)$, or $(/\wedge)$, then \mathbf{A} is integral, and hence is a Boolean algebra.*

Proof. Corollary 4.3 settles the claim if \mathbf{A} satisfies any of $(\cdot\wedge)$, $(\wedge\cdot)$, $(\wedge\backslash)$, or $(/\wedge)$. We therefore prove the claim for \mathbf{A} satisfying $(\backslash\vee)$; it will follow if \mathbf{A} satisfies $(\vee/)$ by a symmetric argument. Suppose that \mathbf{A} satisfies $(\backslash\vee)$, and note that $e \leq \top$ implies $\top\backslash e' \leq e\backslash e' = e'$. By $(\backslash\vee)$, we have:

$$\begin{aligned} \top &= \top\backslash\top \\ &= \top\backslash(e\vee e') \\ &= \top\backslash e\vee\top\backslash e' \\ &\leq \top\backslash e\vee e'. \end{aligned}$$

Hence $\top\backslash e\vee e' = \top$, so $(\top\backslash e)' \wedge e = \perp$. Because \wedge has a residual \rightarrow in any Boolean residuated binar, we get $e \leq (\top\backslash e)' \rightarrow \perp = (\top\backslash e)'' = \top\backslash e$. By residuating with respect to \cdot , we obtain that $\top \leq e$, and hence $\top = e$. \square

Corollary 4.5. *In a unital Boolean residuated binar each of the identities $(\cdot\wedge)$, $(\wedge\cdot)$, $(\backslash\vee)$, $(\vee/)$, $(\wedge\backslash)$, and $(/\wedge)$ is logically equivalent to the other five.*

The two prelinearity equations (lp) and (rp) are not expressible in the absence of a multiplicative identity e , but for unital residuated binars they enjoy a connection to the nontrivial distributive laws even in the absence of associativity. In particular, inspection of the proofs offered in [2] verifies that in a unital residuated binar satisfying

$$(x\vee y)\wedge e = (x\wedge e)\vee(y\wedge e),$$

each of $(/\wedge)$ and $(\vee/)$ implies (lp) , and each of $(\wedge\backslash)$ and $(\backslash\vee)$ implies (rp) . Without associativity, the converse implications fail. To see this, we may define a five-element residuated binar \mathbf{A}_7 whose lattice reduct is pictured in Figure 4. The multiplication \cdot on \mathbf{A}_7 is given in the following table:

\cdot	\perp	a	b	e	\top
\perp	\perp	\perp	\perp	\perp	\perp
a	\perp	a	\perp	a	e
b	\perp	\perp	b	b	\top
e	\perp	a	b	e	\top
\top	\perp	a	\top	\top	\top

The residuals \backslash and $/$ are determined uniquely by the above table as well, and with these operations we have $\mathbf{A}_7 \models (lp)$, (rp) , but each of $(/\wedge)$, $(\vee/)$, $(\wedge\backslash)$, and $(\backslash\vee)$ fail in \mathbf{A}_7 . Note also that $\mathbf{A}_7 \not\models (\cdot\wedge), (\wedge\cdot)$, whence prelinearity does not entail either of the latter distributive laws.

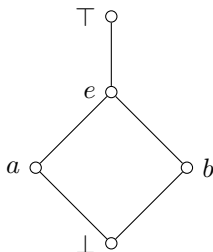


FIGURE 4. Hasse diagram for the lattice reduct of \mathbf{A}_7 .

5. Open problems

Lattice distributivity is a key ingredient in the known proofs of Theorem 2.3, whether purely algebraic or by equivalent frame conditions. We do not know whether any of the implications announced hold in all residuated binars (without assuming lattice distributivity), nor do we know whether any of these implications fail in this more general setting.

When present, a multiplicative identity element plays a decisive role in shaping the connection between the nontrivial distributive laws. Known characterizations of when a residuated binar may be embedded in a unital residuated binar crucially involve terms of the form $x \setminus x$ and x / x (see [1, 6]), and we conjecture that conditions involving terms of this form may provide a more satisfying account of the role of a multiplicative identity in this context. In particular, it would be interesting to identify analogues of prelinearity in the non-unital setting and explicate their connection to the nontrivial distributive laws and semilinearity.

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